

Newton's method

Given a poly f & an approx soln (i.e. $f(a) \approx 0$)
 can "correct" it by solving

$$f(a+\varepsilon) \approx f(a) + \varepsilon f'(a) = 0$$

$$\varepsilon = -\frac{f(a)}{f'(a)} \quad a' = a - \frac{f(a)}{f'(a)}$$

Algebraically: if we have $\underset{\text{poly eqn}}{a}$ " $f=0$ " over
 a ring A and an approx soln $a \in A$ in the
 sense that $f(a)^2=0$ and $f'(a) \in A^\times$

then in fact $a' = a - \frac{f(a)}{f'(a)}$ is a root in A !

(Note: if f poly, $\varepsilon^2=0$ $f(a+\varepsilon) = f(a) + \varepsilon f'(a)$)

(imagination $f' \neq 0$ is a proxy for smoothness)

Reminder S/R is formally smooth if $\mathfrak{I} \otimes A$

and $\mathfrak{I} \subset A$, $\mathfrak{I}^2=0$ \exists lifts in

$$\begin{array}{ccc} R & \xrightarrow{\quad} & A \\ \downarrow & \nearrow & \downarrow \\ S & \longrightarrow & A/\mathfrak{I} \\ x & \longmapsto & \bar{a} \end{array}$$

$$\left[S = R[x]/\mathfrak{I} \right]$$

Philosophically:

smooth = new - problem
- / Newton's method
 $f' \neq 0$

$$f(\vec{a}) = 0$$

$$\frac{\partial f}{\partial a} = 0$$

$$f(a) \in I \quad I^2 = 0$$

i.e. we have noticed that if

$S = R[x]/f$ and roots of f have
the prop. that $f'(a) \in R^*$
then smooth.

S/R

$\frac{k[x]}{f}$ f no repeated
roots
 \Downarrow smooth/ k .

$$\vec{f}(\vec{a} + \vec{\varepsilon}) = \vec{f}(\vec{a}) + (D\vec{f})(\vec{\varepsilon})$$
$$\left(\frac{\partial f_i}{\partial x_j} \right)$$

all entries of
 $\vec{\varepsilon}$ are $D=0$

Already to: (formal statement later)

if "some Jacobian criterion" then formal smoothness

i.e. $S = R[\mathbb{F}_p^2]/I$ w/ some conditions.

Sketch: smoothness is naturally expressed
in terms of poly by presenters.

Def (A, I) is a Henselian pair (A ny $I \triangleleft R$)

if for any poly f and $a \in A$ w/ $f(a) \in I$
 $f'(a) \in A^*$ then $\exists a' \in A$ w/ $f(a') = 0$ and

$$a' + I = a + I.$$

Rem: if $I^2 = 0$ then (R, I) is a Hens. pair.

"Prop" w/ some reasonable hypotheses, S/R ^{form.} smooth



$$\begin{array}{ccc} R & \rightarrow & A \\ & \downarrow & \downarrow \\ S & \rightarrow & A/I \end{array}$$

lifts for any (A, I) Hens.

and (A, I) Hens. pr \Leftrightarrow if S/R smooth

\exists lifts

$$\begin{array}{ccc} R & \xrightarrow{\quad} & A \\ \downarrow & \lrcorner & \downarrow \\ S & \xrightarrow{\quad} & A/I \end{array}$$

Def A local ring (R, m) is Henselian if R, m
B a Hens part.

Silly examples

\mathbb{Q} smooth or \mathbb{Q}

$$S = R = \frac{\mathbb{Q}[x]}{x = f}$$

$$A = \mathbb{Q}[\varepsilon]/\varepsilon^2$$

$$x = \varepsilon \rightsquigarrow x = 0$$

$$f(\varepsilon)^2 = 0$$



$$S = \frac{\mathbb{Q}[x]}{x^2}$$

$$A = \mathbb{Q}[\varepsilon]/\varepsilon^3$$

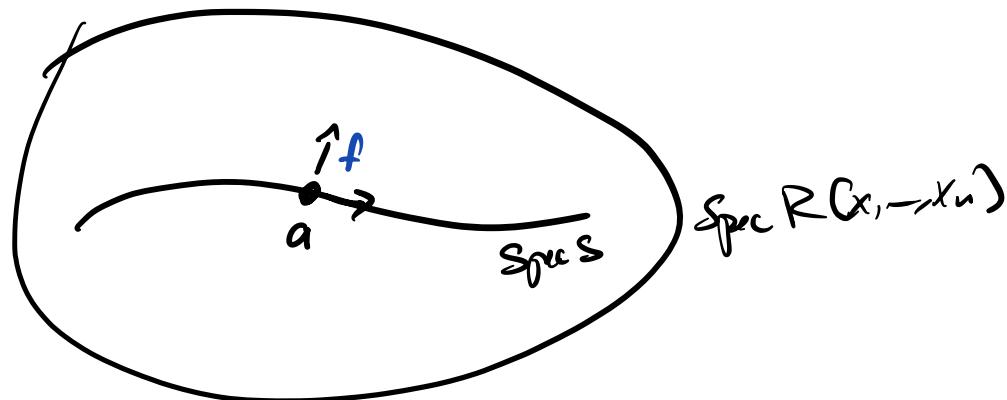
S/\mathbb{Q} not smooth.



Geometric sense of smoothness?

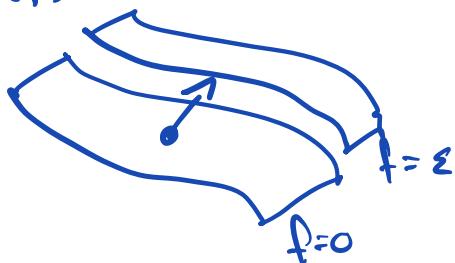
Idea: S/R smooth $S = R[x_1, \dots, x_n] / J$

where near every point (a_1, \dots, a_n)
can find generators (near \vec{a}) f_1, \dots, f_r
w/ $\dim S = \dim R + n - r$
and where f_1, \dots, f_r act like cords
at p



key to cartesian interpretation

$f \in Z(f)$



2 interpretations
of indep. func. cut off right #.

- regular sequences
- differentials

$$\text{if } X \subset A_k^n \quad a \in A_k^n(k) \\ a \in X(k)$$

makes sense to ask if X cut out at a
by a reg. sequence.

$$\mathcal{O}_{X,a} = \mathcal{O}_{A_k^n(a)} / (f_1, \dots, f_r)$$

f_1, \dots, f_r indep in $M_{\alpha}^{(n)}/M_{\alpha}^2$

Can ask for $\mathcal{O}_{X,a}$ my local
all at X scheme + metric.

Def of X regular scheme \nearrow

(Condition: $X \rightarrow Y$ smooth, Y sm $\Rightarrow X$ regular)

Def fix \mathbf{y} smooth if f is loc. of finite presentation

by $\mathrm{Spec} S \rightarrow \mathrm{Spec} R$ S/R fm-smooth
by extension.

Warning: if X/k field regular, $X \rightarrow \mathrm{Spec} k$ need not be smooth.

ex: $k = \mathrm{char} p$ $\frac{k[x]}{x^p - a}$ is a field. \Rightarrow regular.
 $a \notin k^p$

but $L = \frac{k[x]}{x^p - a}$ not smooth.
 $\bar{x} = \alpha \in L$ satisfies $\alpha^p = a$

Claim: formal that S/R smooth T/R alg

$\Rightarrow S \otimes_R T/T$ smooth.

$$\text{note } T = S = k[x]/x^p - a = L$$

$$\begin{aligned} S \otimes_R T &= L \otimes_k L = \frac{L[x]}{x^p - a} \\ &= \frac{L[x]}{x^p - \alpha^p} \\ &= L[\bar{x}]/(\bar{x} - \alpha)^p \end{aligned}$$

$$x - \alpha = \varepsilon$$

$$= L[\varepsilon]/\varepsilon^p$$

$$\underline{\mathcal{E}_{X'}} \left(k[x, \varepsilon]/\varepsilon^2 \right) / \left(k[\varepsilon]/\varepsilon^2 \right) \text{ smooth}$$

$$k[x]/k \text{ smooth}$$

Def Given S/R , define $\mathcal{R}_{S/R}$

free S -module w/ generators $da, a \in S$ modulo
the submod gen. by

- $d(a+b) = d(a) + d(b) \quad a, b \in S$

- $d(ab) = ad(b) + b d(a)$

- $d_r = 0 \quad r \in R$

$\mathcal{R}_{S/R}$ is a repository for differentiation

$$\begin{array}{ccc} S & \xrightarrow{d} & \mathcal{R}_{S/R} \\ a & \longrightarrow & da \end{array}$$

Def A derivation $S \xrightarrow{F} M$ (M an S -mod)

is a map which is an R -lin map i;
 $\varphi(ab) = a\varphi(b) + b\varphi(a)$

$$\mathrm{Hom}_{S\text{-mod}}(\Omega_{S/R}, M) = \mathrm{Der}_R(S, M)$$

Given $B = A \oplus M$ M an A -module
 w/ w stroke given by $(a+m)(b+n)$

$$\underbrace{ab}_{A} + \underbrace{an+bm}_{M}$$

$I = M$ square o id

$$S \xrightarrow{\psi} A \quad R\text{-alg.}$$

$$\tilde{\psi} \rightarrow A \oplus M$$

$$\tilde{\psi}(x) = \psi(x) + \psi'(x)$$

$$\psi': S \rightarrow M$$

$$\tilde{\psi}(xy) = \psi(xy) + \psi'(xy)$$

$$\tilde{\psi}(x)\tilde{\psi}(y) = (\psi(x) + \psi'(x))(\psi(y) + \psi'(y))$$

$$\psi'(xy) = \psi(x)\psi'(y) + \psi(y)\psi'(x)$$

ψ' is a derivation $S \rightarrow M$ wrt o
 S -mod structure on M via

$$S \xrightarrow{\psi} A \otimes M$$

$$\begin{aligned} \left\{ \text{Extensions of } \psi \right\} &= \text{Der}_{\psi}(S, M) \\ &= \text{Hom}_{S\text{-mod}}(SL_{A \otimes M}, M) \\ &\quad (\text{via } \psi) \end{aligned}$$

Def S/R ^{formally} \hookrightarrow smooth if given $I \trianglelefteq A$ $I^2 = 0$

and a diagram $\begin{array}{ccc} R & \rightarrow & A \\ \downarrow & & \downarrow \\ S & \longrightarrow & A/I \end{array}$ then $f: S \rightarrow A$

s.t. $\begin{array}{ccc} R & \rightarrow & A \\ \downarrow & \nearrow & \downarrow \\ S & \longrightarrow & A/I \end{array}$ commutes

Def S/R ^{formally} étale if given $I \trianglelefteq A$ $I^2 = 0$

and a diagram $\begin{array}{ccc} R & \rightarrow & A \\ \downarrow & & \downarrow \\ S & \longrightarrow & A/I \end{array}$ then $f: S \rightarrow A$

s.t. $\begin{array}{ccc} R & \rightarrow & A \\ \downarrow & \nearrow & \downarrow \\ S & \longrightarrow & A/I \end{array}$ commutes

Def S/R ^{formally} unramified if given $I \triangleleft A$ $I^2 = 0$

and a diagram

$$\begin{array}{ccc} R & \xrightarrow{\quad} & A \\ \downarrow & & \downarrow \\ S & \xrightarrow{\quad} & A/I \end{array}$$

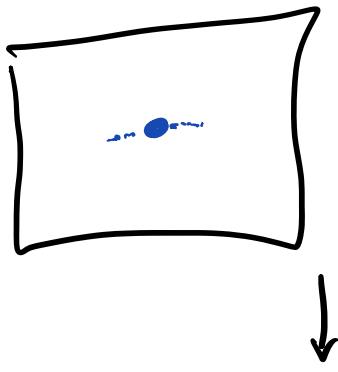
then $f \wedge S \rightarrow A$
at most one

s.t.

$$\begin{array}{ccc} R & \xrightarrow{\quad} & A \\ \downarrow & \nearrow & \downarrow \\ S & \xrightarrow{\quad} & A/I \end{array}$$

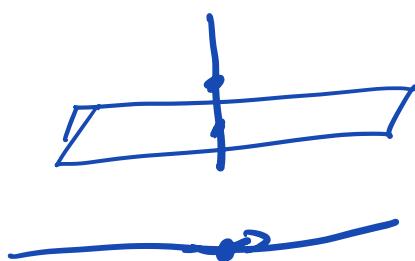
commutes

so: $f_{\text{ét}} = f_{\text{sm}} + f_{\text{unram.}}$



✗ smooth

— étale
supposed to capture
"local differentiation"



Def smooth = f. sm + loc. frk punctures
(on either)

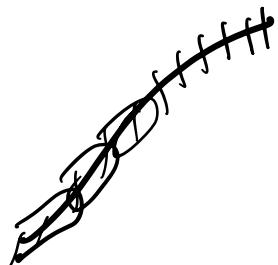
Def étoile = f. ét + loc. frk punctures
(on either)

Def unram = f. unr + loc. frk type
(on either) (f.s. gen alg.)
relations...?

w → x

↓ closed which, ? } f-smooth
z → y

$$\omega_w^2 = 0 \text{ in } \Omega_z$$



$$\begin{array}{ccc} \text{Spec } A_{/\mathcal{I}} \rightarrow X & & \\ \downarrow \text{, } \exists \nearrow & \downarrow \text{f-sm} & \mathcal{I}^2 = 0 \\ \text{Spec } A \rightarrow Y & & \end{array}$$